



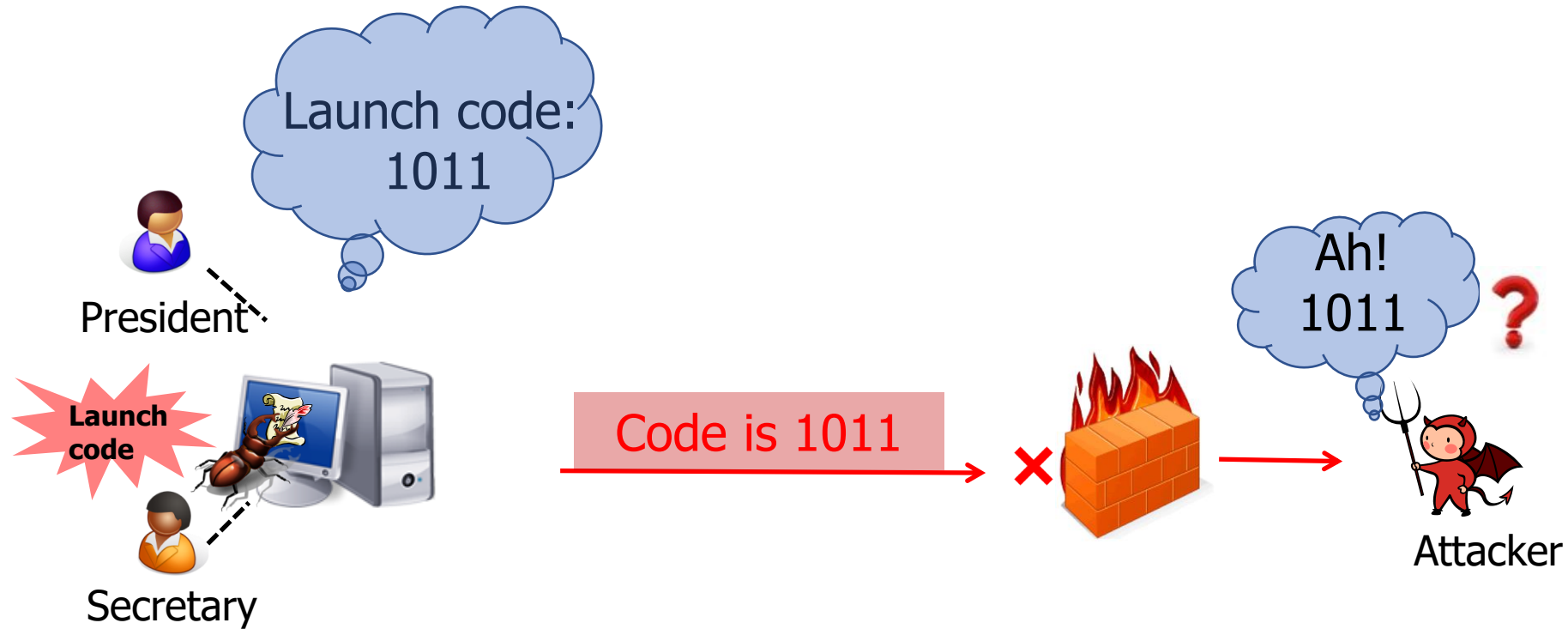
NetWarden: Mitigating Network Covert Channels While Preserving Performance

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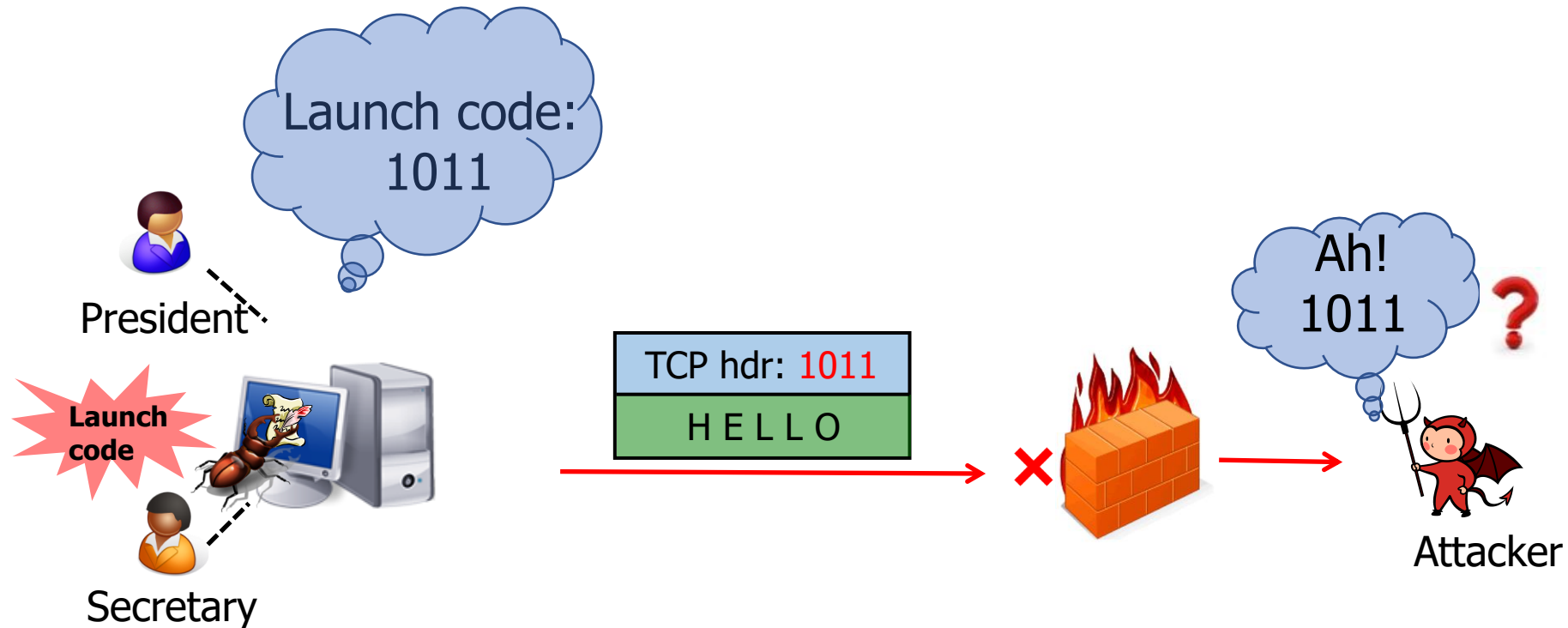
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Motivation: Mitigating network covert channels

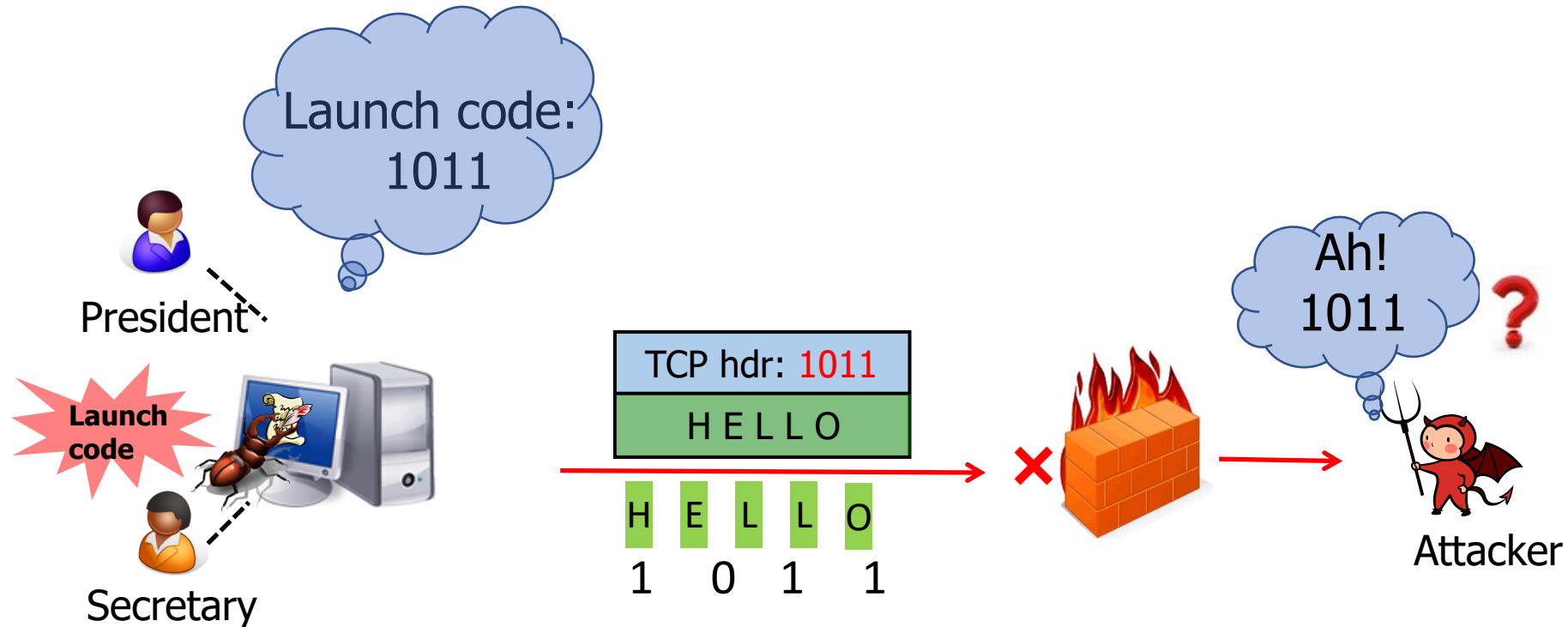


Motivation: Mitigating network covert channels



- Covert channels:
 - Storage channels: changing the packet **header fields**.

Motivation: Mitigating network covert channels



- Covert channels:
 - Storage channels: changing the packet **header fields**.
 - Timing channels: changing the **timing** of packets.

Covert channels are a longstanding problem

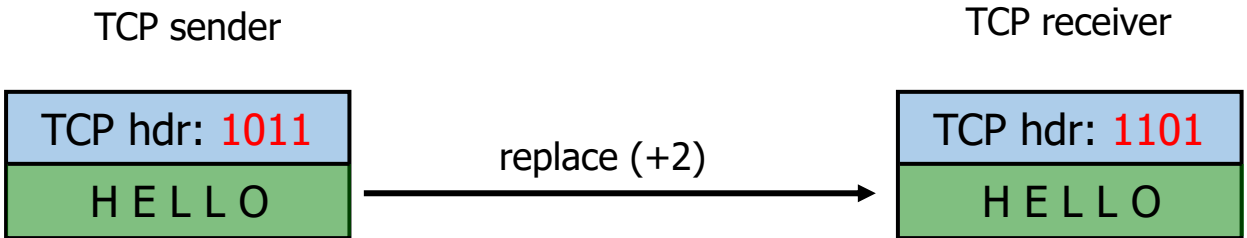
A note on the confinement problem

Author:  [Butler W. Lampson](#) [Authors Info & Affiliations](#)

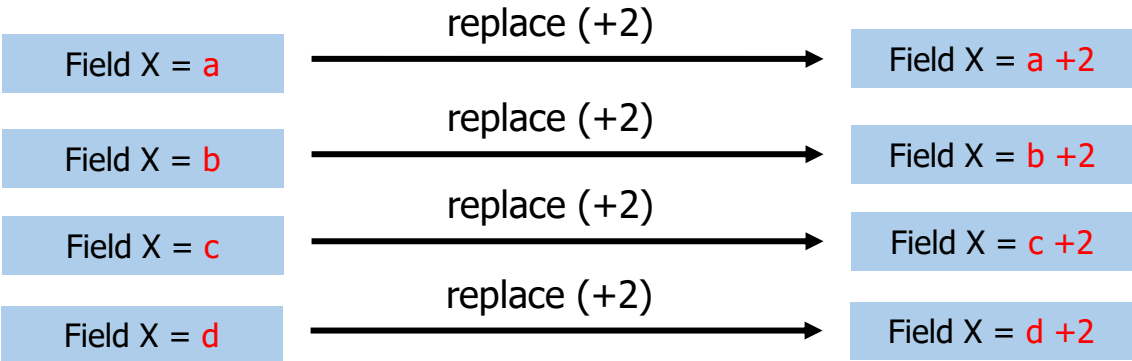
Publication: Communications of the ACM • October 1973 • <https://doi.org/10.>

- They can leak data over long distance effectively
 - Covert storage channels
 - TCP ISN (1997), TTL (2004), Partial ACK (2009)
 - Covert timing channels
 - IP-layer (2004), TCP-layer (2008), PHY-layer (2014)
- Major security standards require protection against them
 - E.g., Common Criteria

State of the art: Storage channel defense



Repeat for **EVERY** packets for **Tbps** traffic

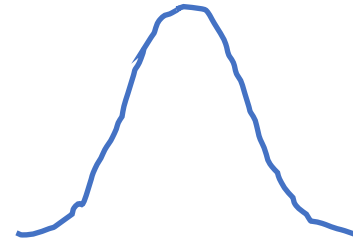


...

- State-of-the-art solutions are software-based
 - Detection: Per-packet header inspection
 - Mitigation: Per-packet header modification
- As a result, they are very **inefficient!**

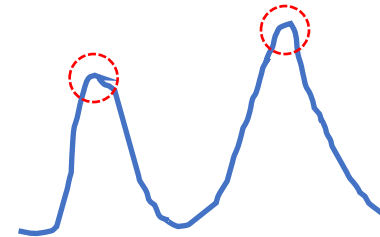
State of the art: Timing channel detection

Normal traffic:



One peak

With channel:

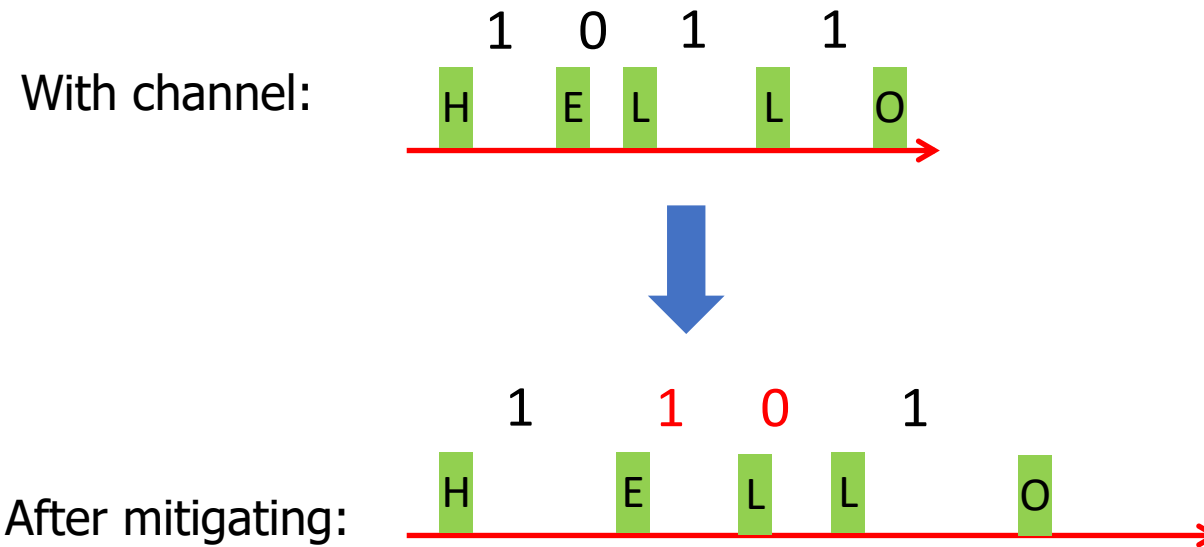


Two peaks

Distribution of packet gaps

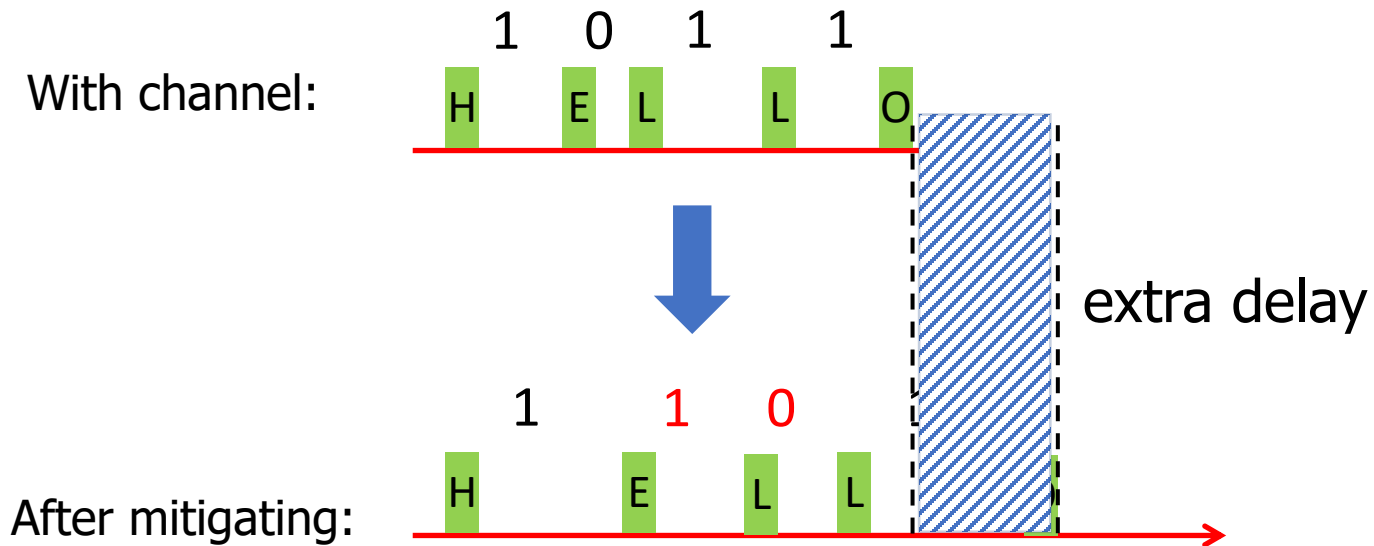
- Detection: Statistics-based tests over packet gaps
 - Looking for signs of statistical deviation
 - → Not always accurate

State of the art: Timing channel mitigation



- Mitigation: Add random delay to each packet
 - Destroy the original timing of the packets

State of the art: Timing channel mitigation



- Mitigation: Add random delay to each packet
 - Destroy the original timing of the packets
- It will **increase the latency** of TCP connections

Problem: Performance penalty

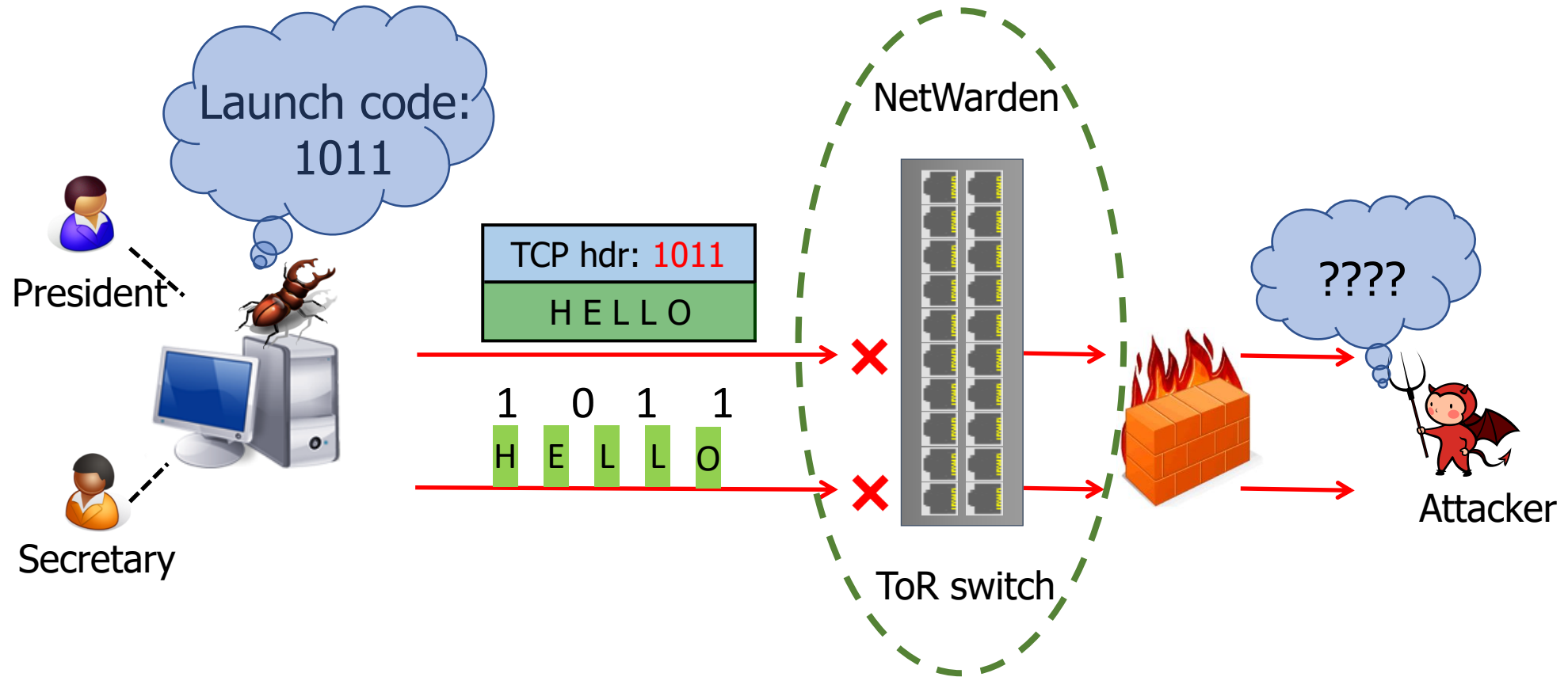
- Detection:
 - **Per-packet** inspection required
 - **Software** cannot keep up with Tbps traffic
- Mitigation:
 - **Adding random delay to each packet** → Increase latency
 - **Collateral damage** → Affects legitimate traffic (e.g., false positives)

Key question

Can we mitigate covert channels while **preserving performance?**



Approach: NetWarden



- **NetWarden:** A performance-preserving covert channel defense.

Key challenges and solutions

- **Key principle:** Hardware/software co-design
- **Challenge #1:** Efficient detection
 - **Solution:** Using **P4** switches
- **Challenge #2:** Performance-preserving mitigation
 - **Solution:** Performance “boosting”

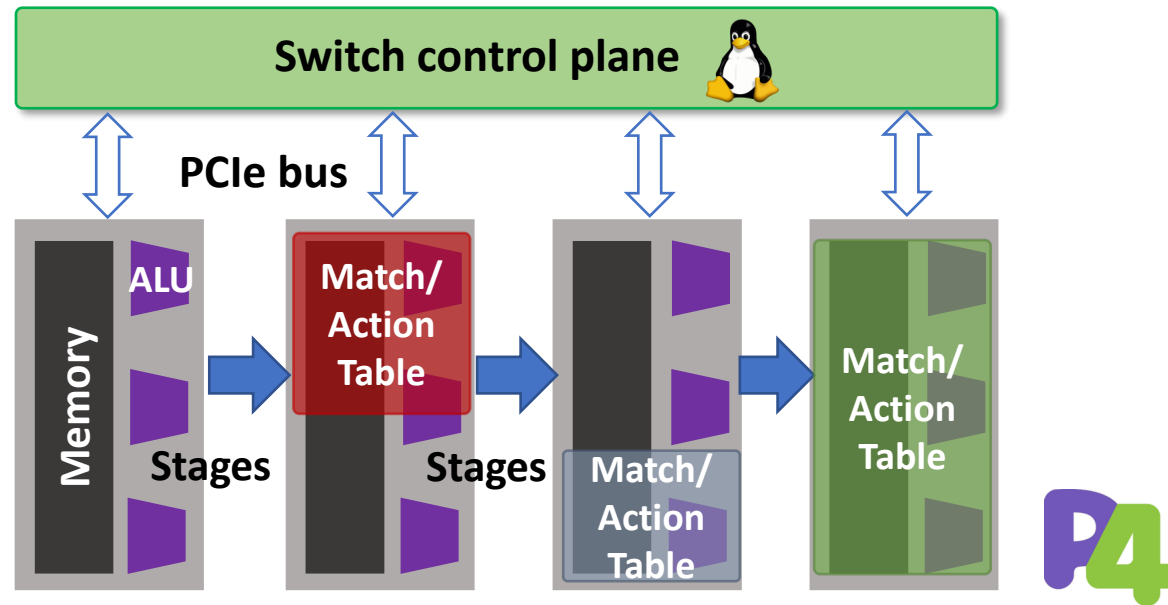
Outline

- ✓ - Motivation: Mitigating network covert channels
- ✓ - State of the art: Performance penalty
- ✓ - Approach: NetWarden
- ➔ - NetWarden design
 - Principles of hardware/software co-design
 - Challenge #1: Efficient detection
 - Challenge #2: Performance-preserving mitigation
- Evaluation
- Conclusion

Key principle: Hardware/software co-design

- A **generic** principle that is applicable to many P4 applications.

P4 switch anatomy

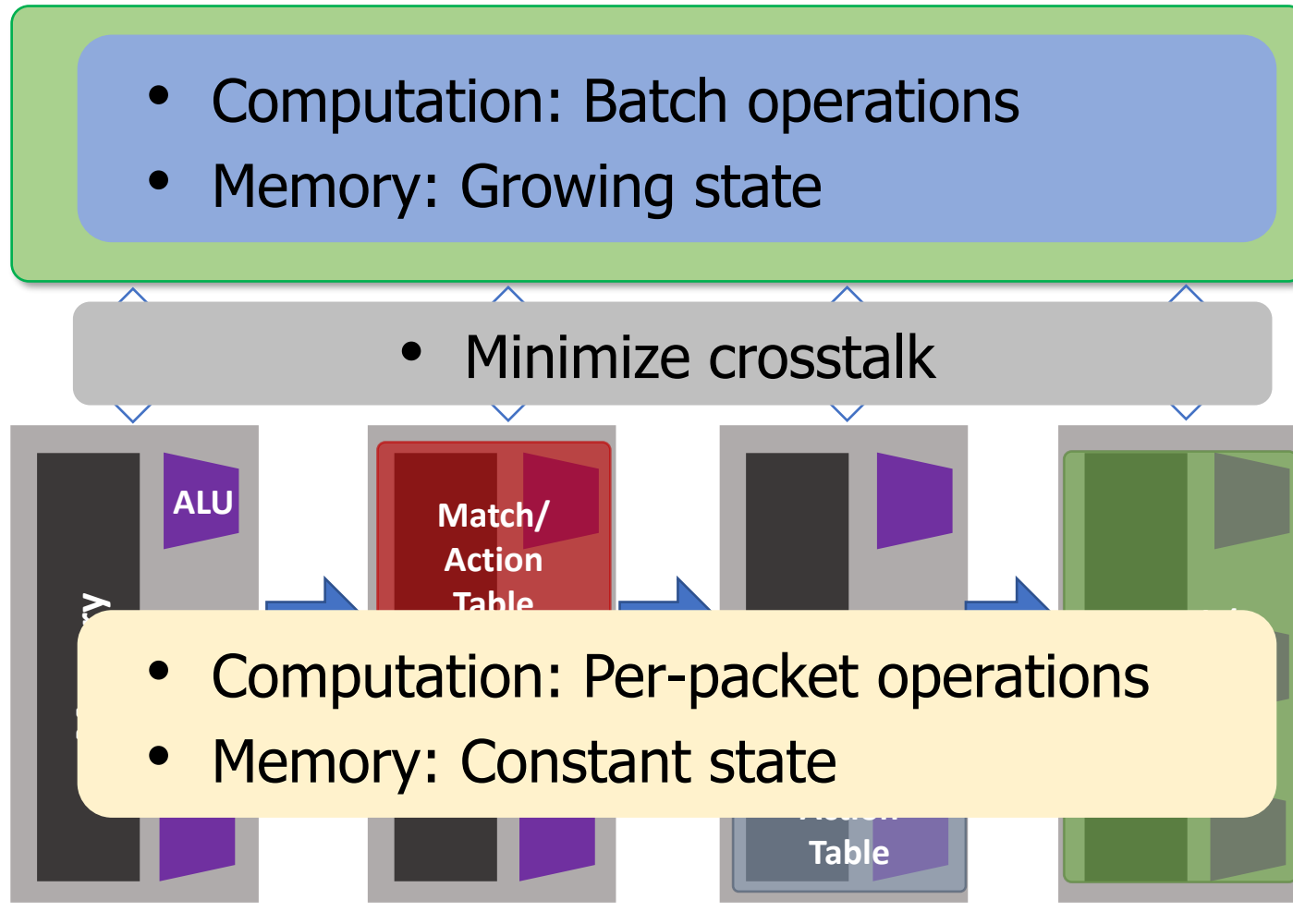


- Data plane
 - Header modification, ns timestamp, per-flow state, line speed
 - Limited memory, simple math computation
- Control plane
 - Abundant memory, complex math computation
 - Software speed



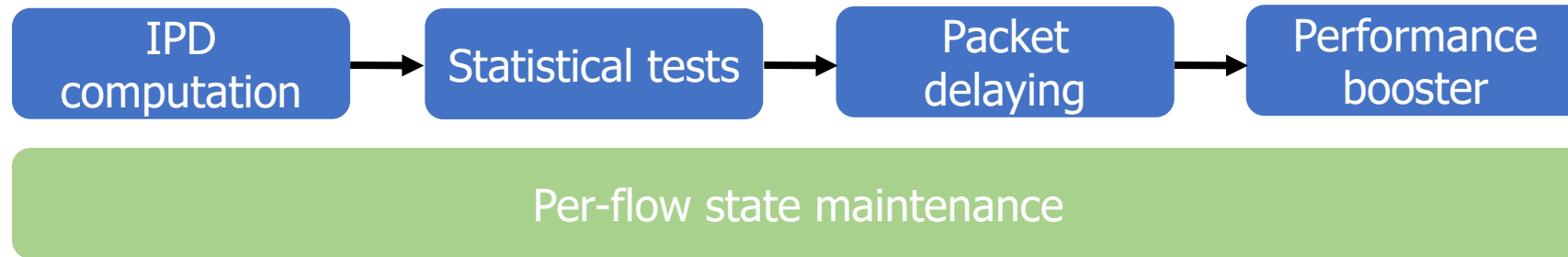
Complementary!

Key principle: Hardware/software co-design

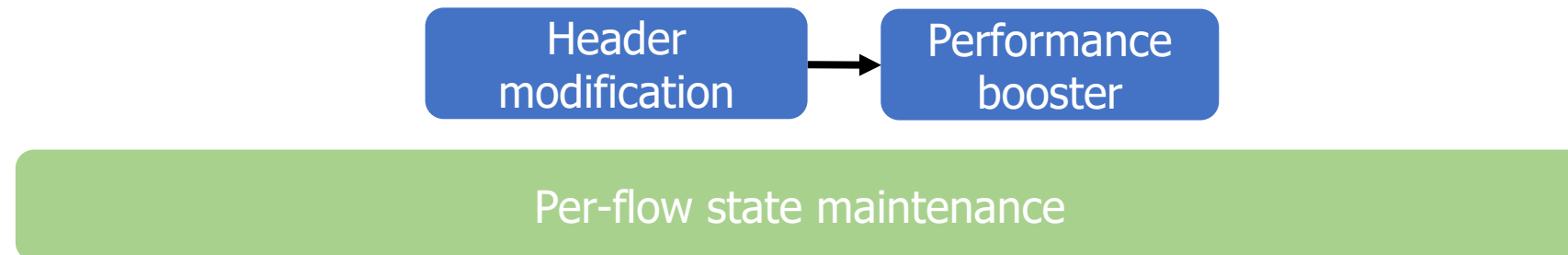


Covert channel defense roadmap

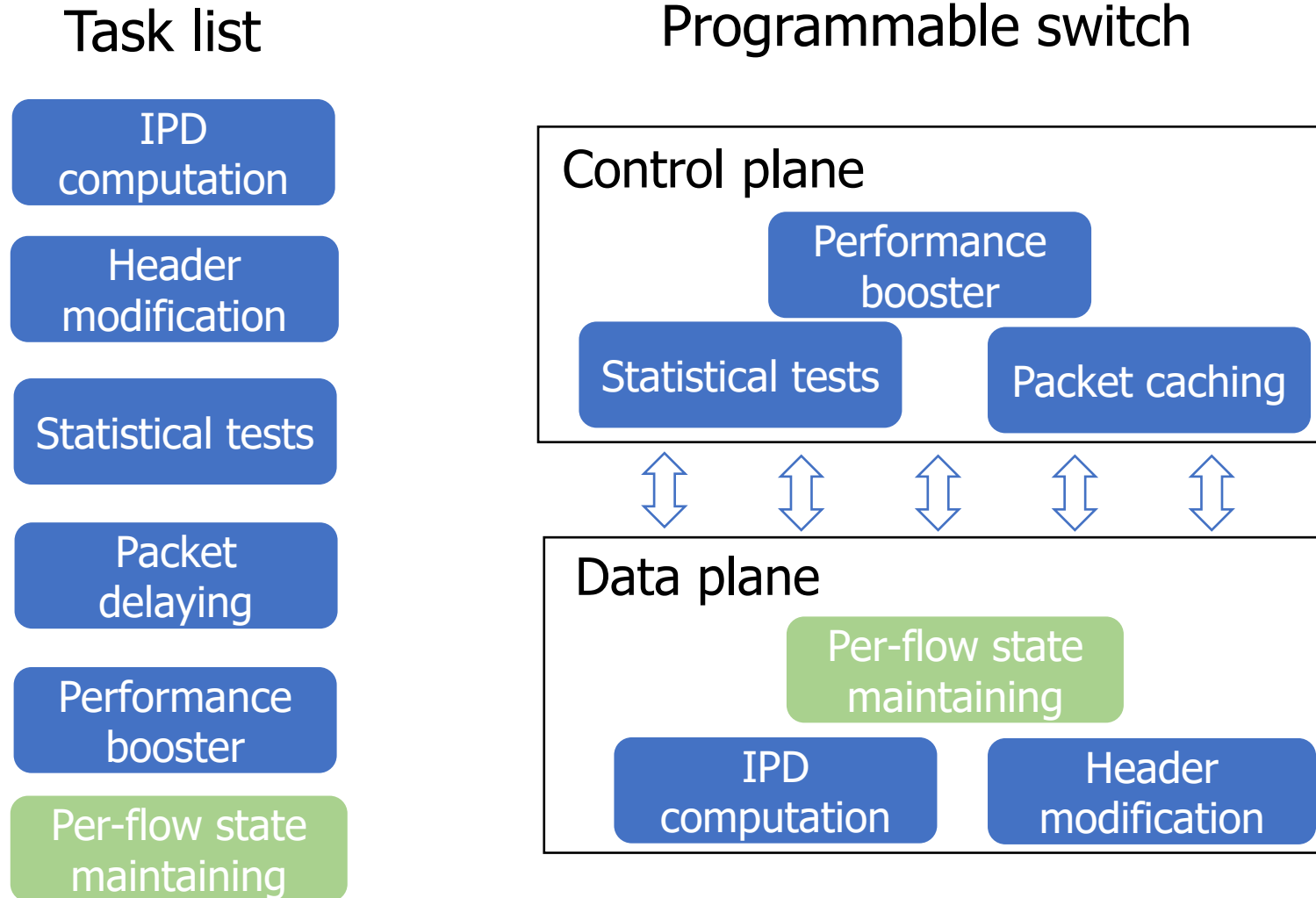
- Time channel defenses



- Storage channel defenses



Applying the hardware/software co-design principle

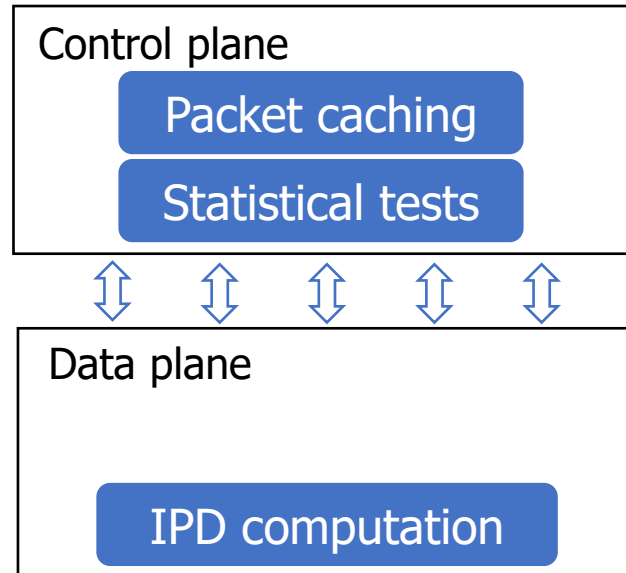


Challenge #1: Efficient detection

- **Solution:** Build efficient detections in **P4 switches** by applying the hardware/software co-design principle.

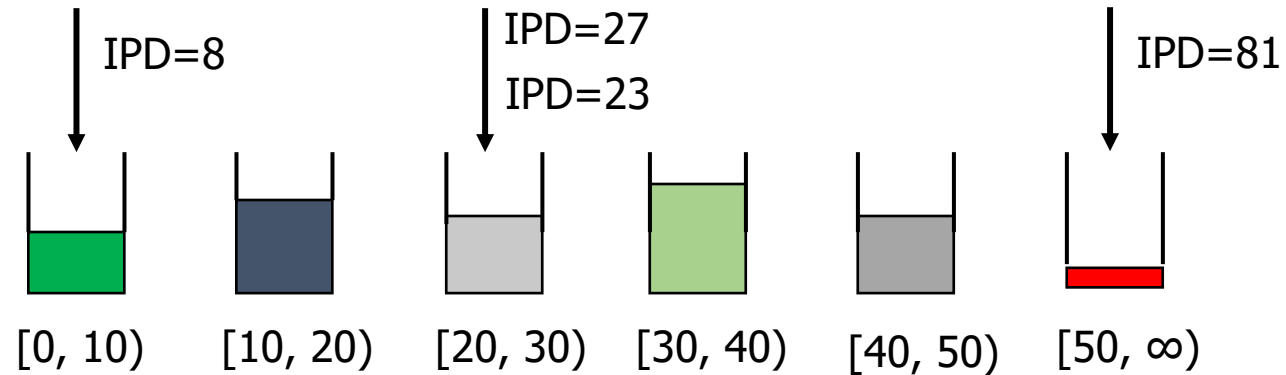


How to compute inter-packet delay (IPD)?



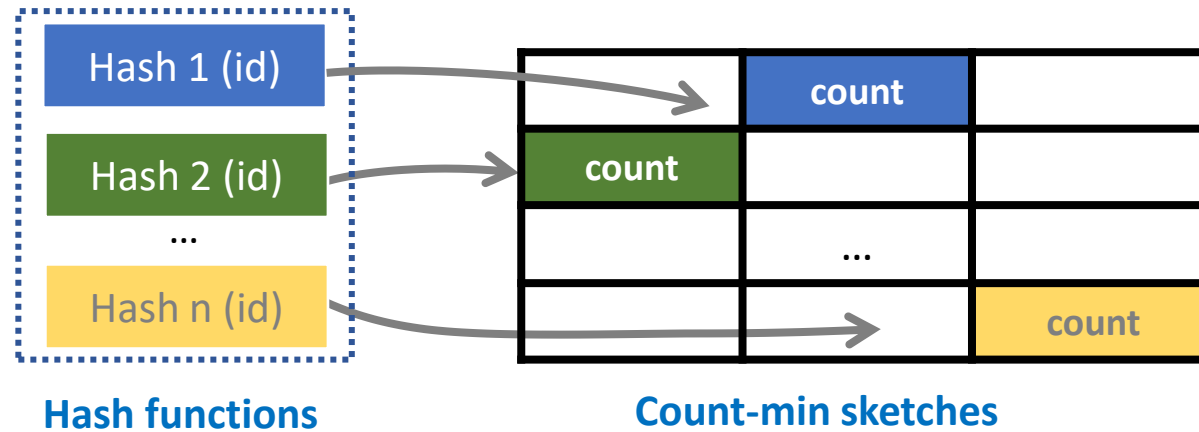
- Solution: Using ns-timescale P4 timestamps
 - $IPD = \text{Current timestamp} - \text{last timestamp}$
- Send all IPDs to the control plane directly?
 - **Crosstalk minimization principle!**

How to store IPDs in a memory-efficient manner?



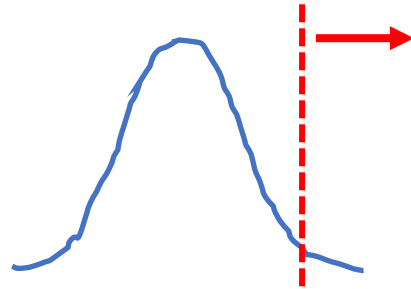
- Solution: IPD intervalization
 - Store IPD interval counters rather than exact IPDs

How to reduce per-flow memory consumption?



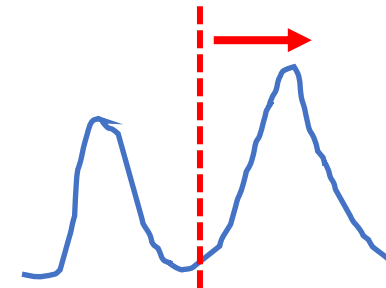
- Solution: IPD sketching
 - Trade off per-flow accuracy for space saving

How to reduce IPDs sent to the control plane?



One peak

Normal traffic pattern



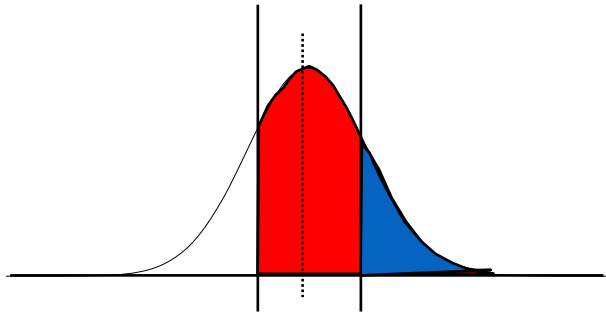
Two peaks

Suspicious pattern

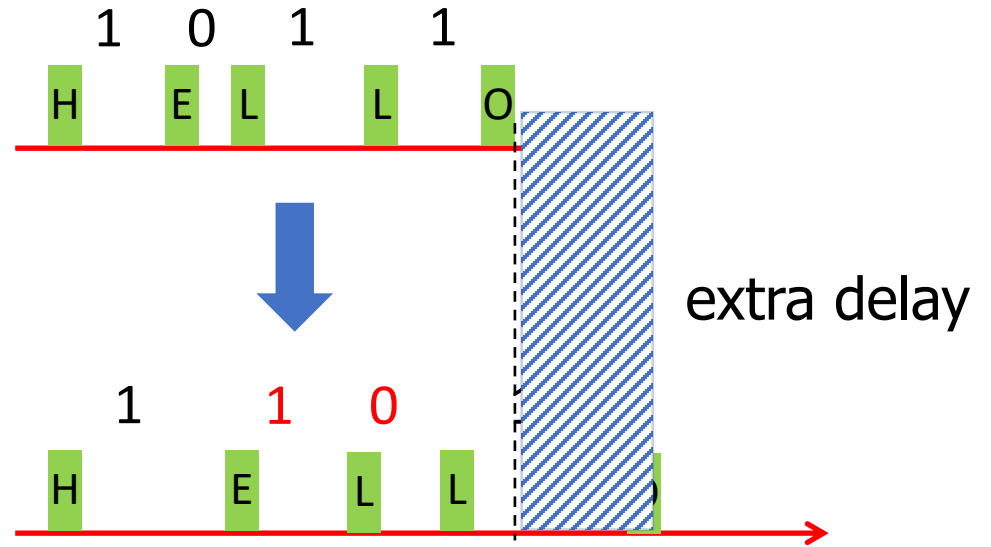
- Solution: IPD pre-check in the data plane
 - Do a quick check and only send suspicious flow IPDs to the control plane.

How to mitigate covert timing channels?

Statistical tests



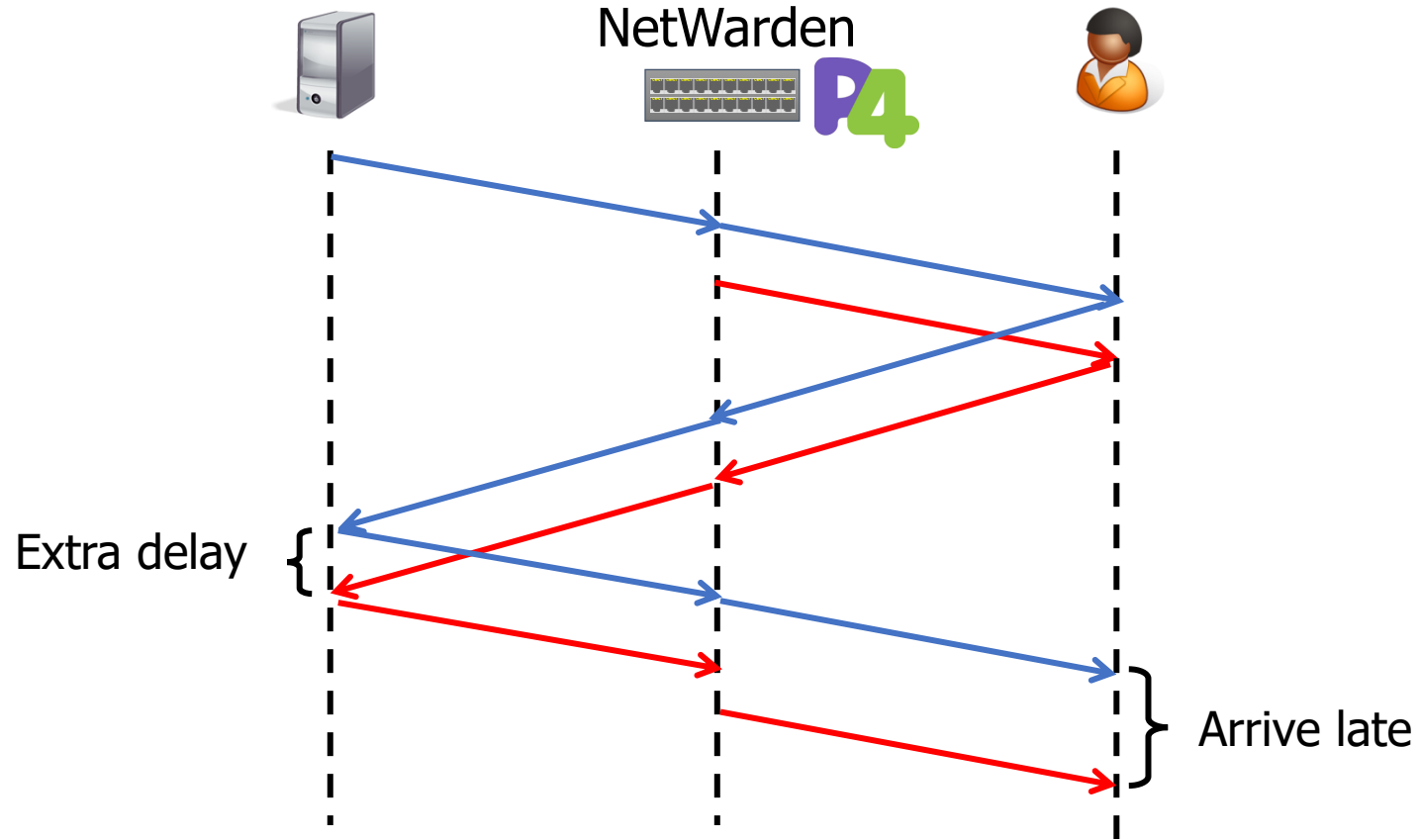
With channel:



After mitigating:

- Solution: Using the Control plane
 - Performs statistical tests
 - Adds random inter-packet delay by caching
- Note: This incurs extra delay.

Challenge #2: Performance-preserving mitigation



- Problem: Existing mitigations incur performance loss.

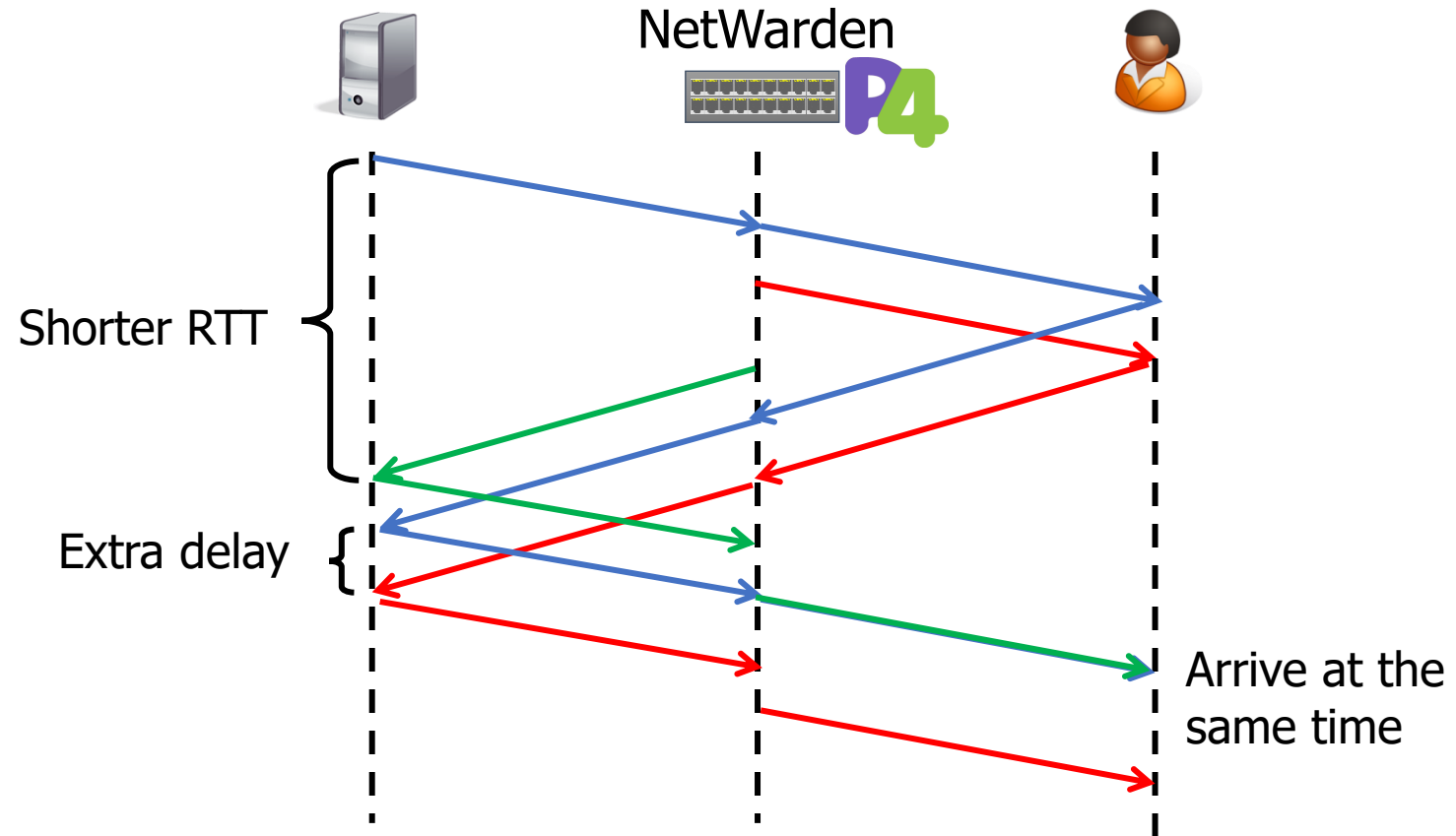
Challenge #2: Performance-preserving mitigation

- Solution: Temporarily boosting TCP performance to **neutralize** the performance penalty.
- Two boosters:
 - **ACK booster**: Generate ACK packets in advance.
 - **Receive window booster**: Enlarge receive window field temporarily.

Challenge #2: Performance-preserving mitigation

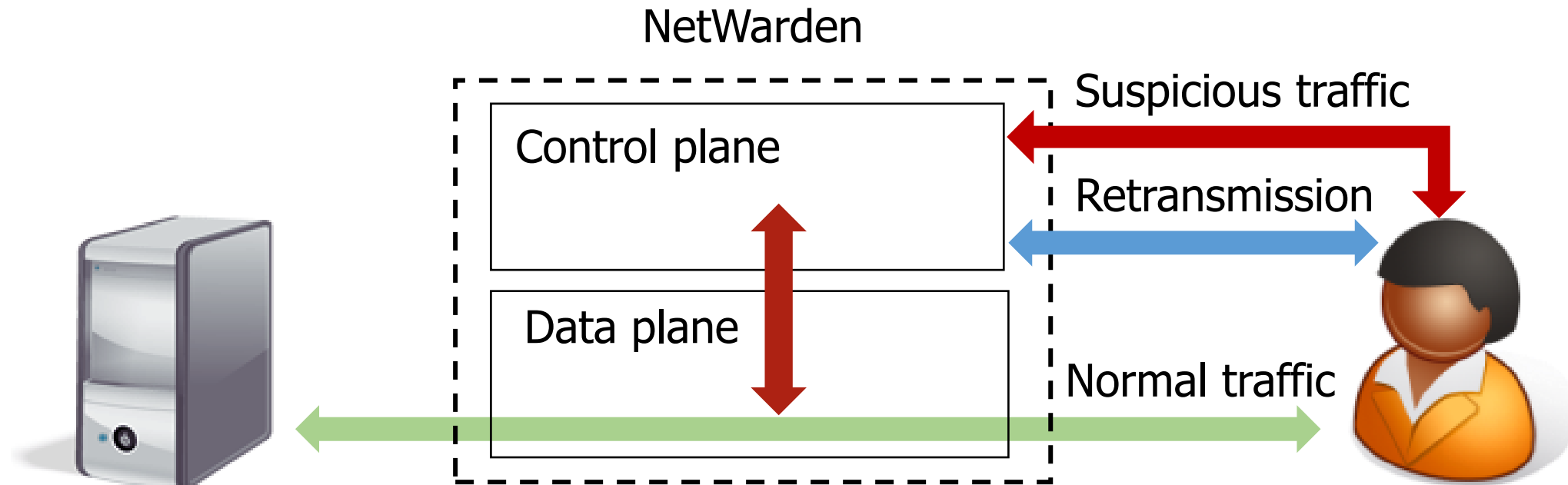
- Solution: Temporarily boosting TCP performance to **neutralize** the performance penalty.
- Two boosters:
 - **ACK booster**: Generate ACK packets in advance.
 - Receive window booster: Enlarge receive window field temporarily.

Boosting performance: ACK booster



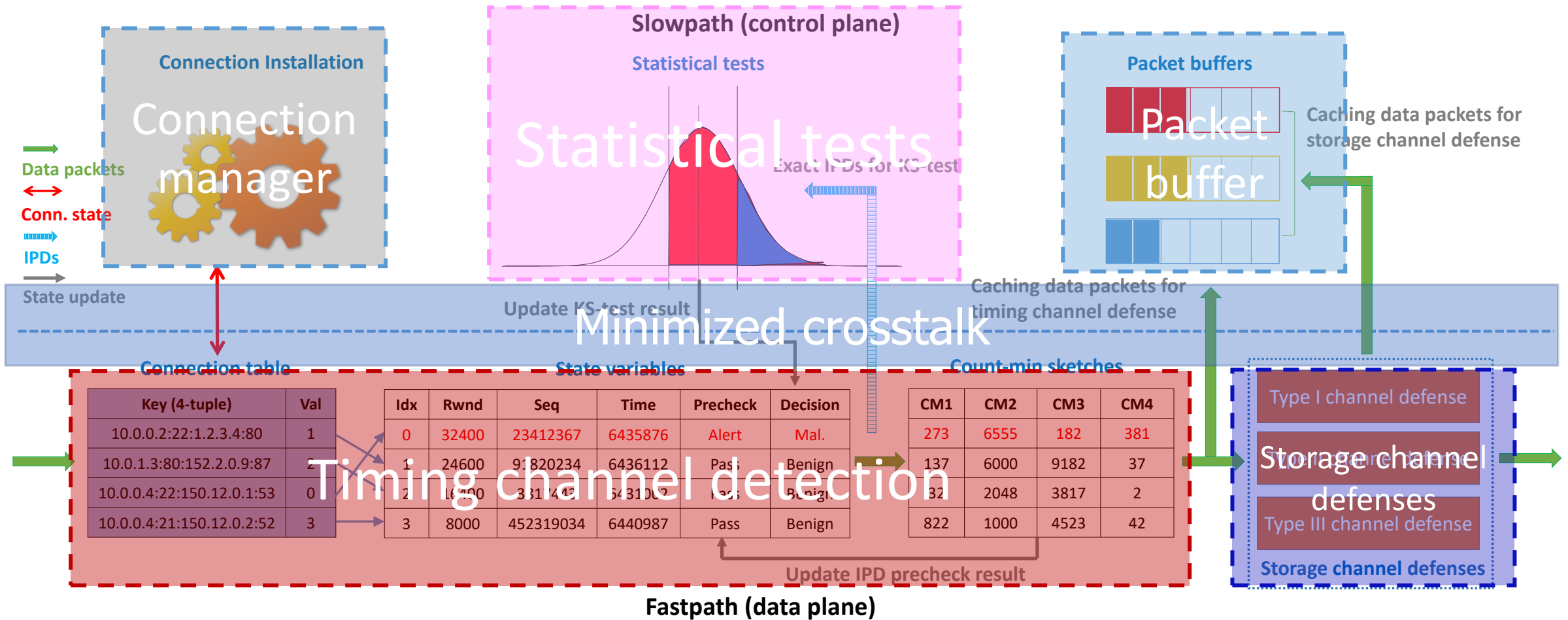
- Creates the **illusion** of a shorter latency as perceived by the sender.

Works as a TCP proxy



- NetWarden works as a **TCP proxy** for malicious traffic.

NetWarden panorama



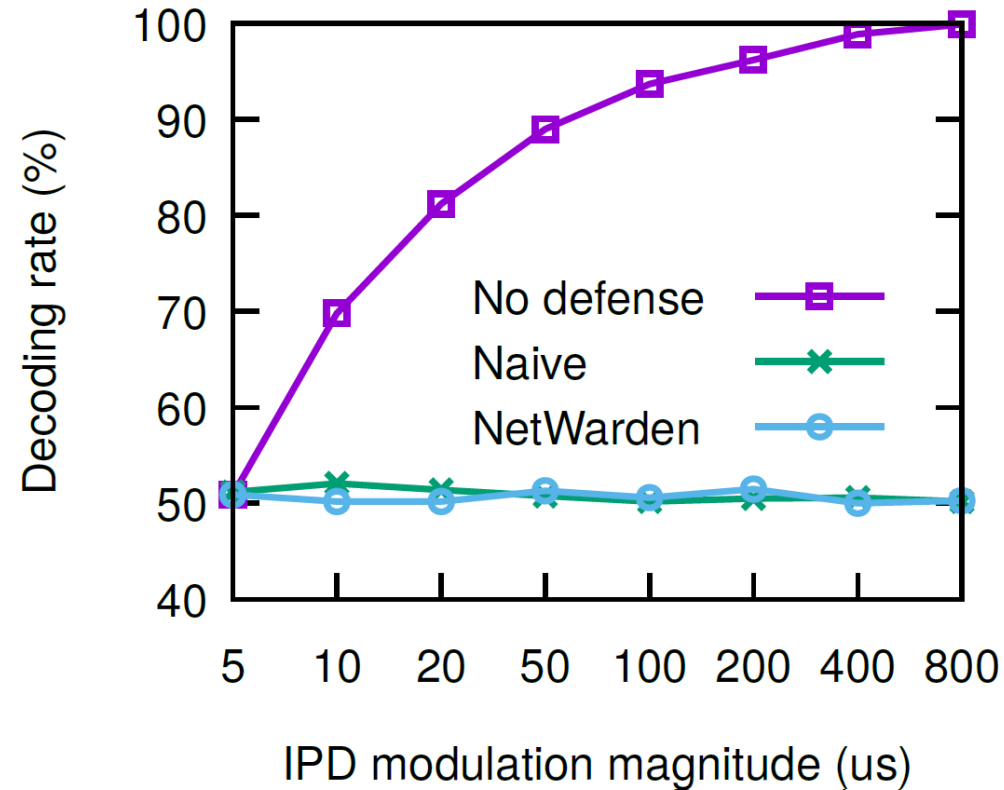
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Evaluation setup

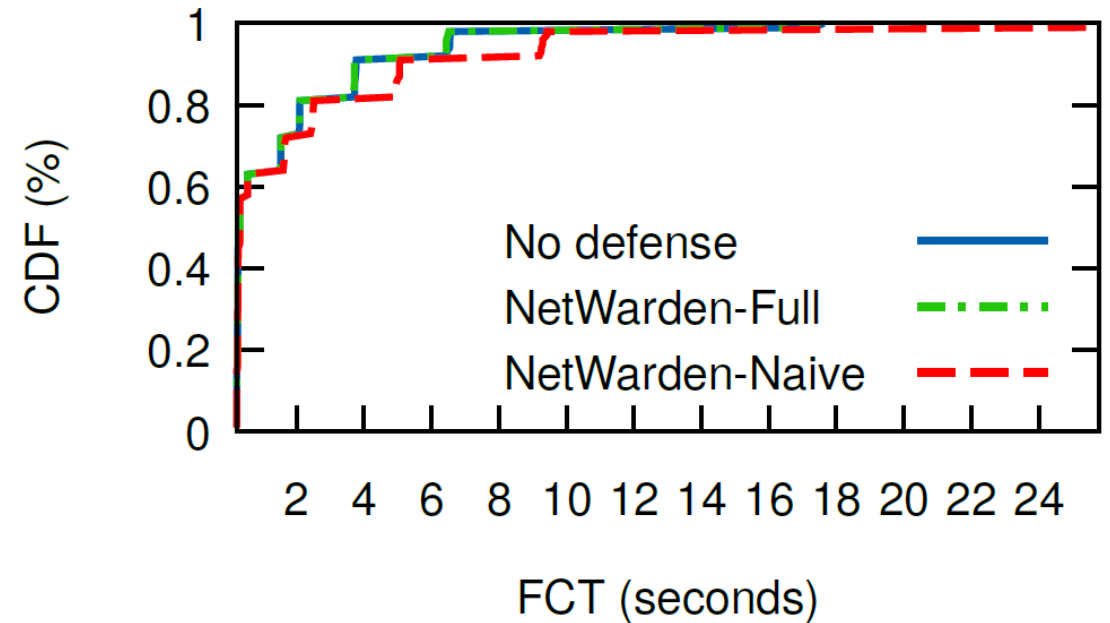
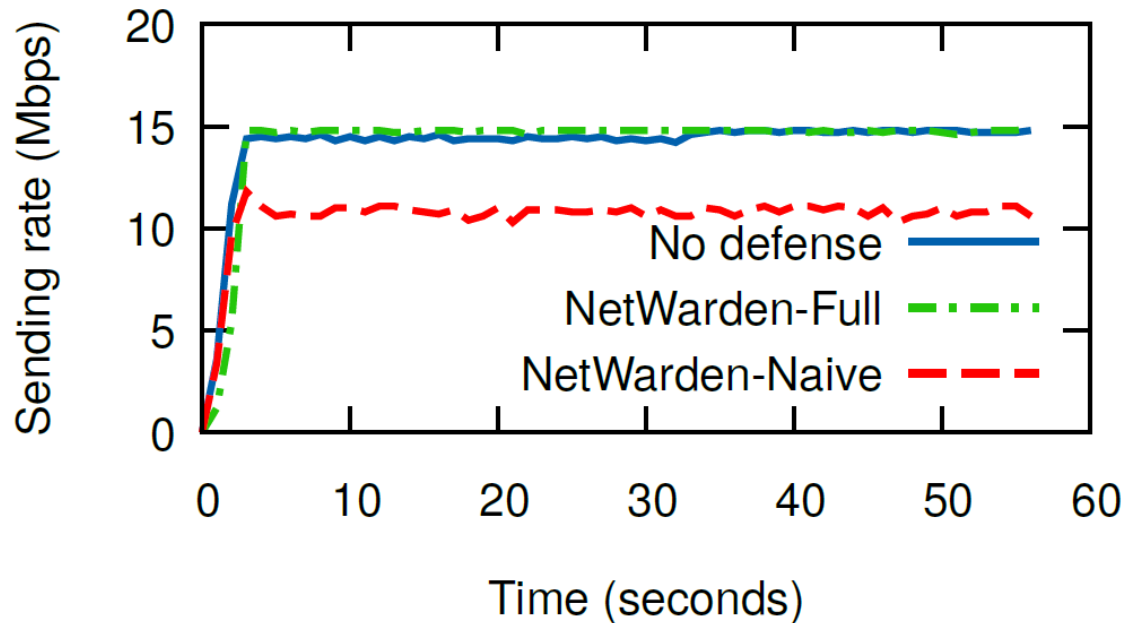
- NetWarden prototype:
 - Runs in Tofino Wedge 100BF-32X switch.
 - 2500 LoC of P4 + 3000 LoC of C+Python
- Threat model:
 - A compromised server + a trusted P4 switch running NetWarden
 - Leak a 2048-bit RSA key via covert channel.
- Workloads:
 - Web search (Alizadeh-SIGCOMM'15)
- Baseline:
 - Defenses without performance boosting

How effective is NetWarden in covert channels mitigation ?



- Naïve defense: renders decoding to a random guess.
- NetWarden: very close to a random guess.
- NetWarden can mitigate covert channels **effectively**.

How well does NetWarden preserve performance?



- Naïve defense incurs 25% performance penalty.
- NetWarden only has 1% performance deviation.
- NetWarden can mitigate covert channels **with minimal performance loss.**

See more evaluation results in our paper

- How effective is NetWarden in covert channels mitigation?
- How well does NetWarden preserve performance?
- How well does NetWarden work with different TCP variant?
- How scalable is NetWarden?
- How much overhead does NetWarden incur?
- How well does NetWarden work with complex applications?
- How robust is NetWarden in self attacks?

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Conclusion

- Motivation: Mitigating network covert channels
- Key limitation of existing approaches:
 - **Performance penalty**
- Our approach: **NetWarden**
 - Principles of hardware/software co-design
 - Efficient detection and mitigation
 - Performance preservation
- Evaluation:
 - **Mitigates covert channels with minimum performance loss!**



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