

TCP-INT: Lightweight INT in TCP Transport

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Agenda

- Motivations
- TCP-INT High-Level Design
- Example use cases
- Implementation (Intel)
- Demo
- Congestion Control Enhancement
- Implementation (Baidu)
- Baidu Vision & Planned Usage
- Roadmap

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Motivations

- The separation between storage and computation requires a low latency, high throughput fabric to realize the benefit
- BaiduRPC over RDMA is used in the storage cluster
- Congestion & bottleneck exist between Storage client and master/block server due to the use of TCP



Motivations (2)

Feedback on INT

- 1. Flow Scale
- 2. Telemetry "for Control"
- 3. Direct mapping to application session/message



TCP-INT High-Level Design (3)



- In-band Network Telemetry (INT) embedded in the TCP header (as TCP Option)
- Correlates fabric telemetry (Q depth) with TCP states (congestion window)
- Lightweight: metrics aggregated (max or sum) over switch hops
- INT enriched with end-host metrics



TCP-INT High-Level Design







TCP-INT High Level Design (2)

New TCP Option – TCP-INT

		Updated by the switches		back to the senders, ignored by the switches			
Option- kind (1B)	Option- length (1B)	INTval (1B) (Scaled summation of utilization or queue depth)	INTecr (1B)	ID (1B)	IDecr (1B)	Lat (3B)	LatEcr (3B)
0x72	0x0C	<pre>if qdepth < qdepth_threshold: Bit [7]: 0 Bits [0-6]: util >> y (saturates if above max) else: Bit [7]: 1 Bits [0-6]: qdepth >> x (saturates if above max)</pre>	INTval Echo Reply	IP.TTL	IP.TTL Echo Reply	Lat += switchLat	Lat Echo Reply

Example Use Cases

Extend Host Linux TCP Toolset with TCP-INT Information

- Quick view into network state
- Compare different congestion control algorithms
- Develop new congestion control algorithms
- Debug QoS configuration in the fabric (e.g. WRED thresholds)





Implementation

Intel Implementation Overview

End-hosts



Switch





Implementation (2)



Host-side eBPF-based implementation

- eBPF callbacks for TCP options (kernel 5.10). TCP-INT eBPF code is called when:
 - new TCP connection is established -> enabling option callbacks for the lifetime of a flow
 - any unknown TCP option is received -> look for and handle TCP-INT
 - TCP adds options to a new outgoing segment -> initialize TCP-INT (INTval=0)
- SK buff local storage (kernel 5.2) used to keep TCP-INT state (echo TCP-INT b2s)
- Perf events and histograms for live data monitoring
 - Exportable via other user-space applications (E.g. gRPC client)
- Configuration maps for enable, disable, disable echo
- User-space application loads, controls, and polls data from eBPF program







Switch-Side TCP-INT

- Control planes determines the mapping between (qdepth, txRate) and INTval
- This increases flexibility and allows complex mapping functions that include division



Demo: End-to-End Performance Bottleneck Identification



FIO NVMe Storage Benchmark

- As number of initiators increases, storage application latency increases
- Identifying the root-cause is a big problem in large distributed systems

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Demo: End-to-End Performance Bottleneck Identification



Enhancing CC Algorithms with TCP-INT



Goal: Simple implementations to allow customers add their secret sauce







ECN+INT Congestion Control



ECN+INT Congestion Control

Switch Side

- INTval conveys absolute qdepth

Sender Side

Introduce α_{int} : $\alpha_{int} = qdepth_{smoothed} / qdepth_{target}$

 $qdepth_{smoothed} \leftarrow (1 - g)^*qdepth_{smoothed} + g^*qdepth$ g - smoothing factor

- Combine α_{DCTCP} and α_{int} : $\alpha = \max(\alpha_{DCTCP}, \alpha_{int})$

This approach ensures interoperability with switches that do not support TCP-INT



ECN+INT CC Demo Topology



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Early Performance Data

6K reqs/s

ECN+INT ECN Longest Queue Encountered in path Avg 👻 in each time bucket Moving Average over previous: 2s -286.4kB -40% 184.2kB 124.2kB 64.2kB-Apr 20, 15:08:48.000 Apr 20, 15:09:03.000 Apr 20, 15:08:33.000 Apr 20, 15:09:18.000 Apr 20, 15:09:33.000 Max Link Utilization (%) Max 👗 in each time bucket Moving Average over previous: No Smoothing 🔻 Link util: 100% 100% Apr 20, 15:08:33.000 Apr 20, 15:08:48.000 **157 us average latency** 67 us average latency -43% 126 us p99 tail latency **172 us p99 tail latency** -23%

2.5x*

15K reqs/s

int

*For workloads and configurations visit www.intel.com/PerformanceIndex. Results may vary.

Baidu Implementation



- Congestion window update
 - Baidu implements HPTCP Algorithm
 - HPTCP is based on HPCC algorithm for RDMA networks, ported to TCP
 - Uval is calculated from INTVal field in TCP-INT

HPCC:<u>https://dl.acm.org/doi/pdf/10.1145/334130</u> 2.3342085

- Uval calculaction
 - If INTval[7]=1, INTval[6:0]=qdepth, then: $Uval = \frac{qdepth}{B*RTT} + \frac{txRate}{B} = \frac{qdepth}{B*RTT} + 1$
 - If INTval[7]=0, INTval[6:0]= Uval << bitshift
- Fast retransmit: snd_cwnd = snd_cwnd / 2
- RTO: snd_cwnd = min(snd_cwnd/2, Winit)

Baidu Vision & Planned Usage

- Intend to be used in Storage Client
- Comparison of various technical solutions is in progress
 - SWIFT
 - HPTCP
 - ...
- Aim to compete with the storage networks from top cloud vendors



Roadmap (Intel)

- Open-Source TCP-INT host-side code
- Delivering network telemetry directly to applications
- Scale out testing with 100s of nodes
- Continue research into congestion control improvement and network mgmt. utilizing TCP-INT telemetry



Baidu Roadmap

- Continue to test HPTCP in Baidu's testbed
- Keep optimizing HPTCP congestion control algorithm
- Compare the performance of HPTCP with SWIFT and RDMA schemes
- Deploy the new transport stack in Storage client to improve the entire storage networks

Notices and Disclaimers

- Performance varies by use, configuration and other factors. Learn more at <u>www.Intel.com/PerformanceIndex</u>.
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